

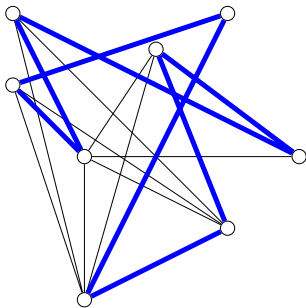
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Hamilton Cycles in Erdős–Rényi  
Subgraphs of Large Graphs

Probabilistic Midwinter meeting  
Umeå Universitet  
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# Hamilton cycles

- A **Hamilton cycle** is a cycle passing through each vertex **exactly once**
- A graph  $G = (V, E)$  is **Hamiltonian** if it contains a Hamilton cycle



## Hamilton cycles

- Minimum degree  $\delta(G) \geq 2$  necessary for Hamilton cycle
- Minimum degree  $\delta(G) \geq n/2$  sufficient for Hamilton cycle (Dirac '52)

In random graph models the gap is usually much smaller

### “Theorem”

*Let  $G$  be a random graph from some well-known model. Then there should be some small  $k$  (usually 2 or 3) such that*

*$G$  is Hamiltonian with high probability*

*if and only if*

*$\delta(G) \geq k$  with high probability.*

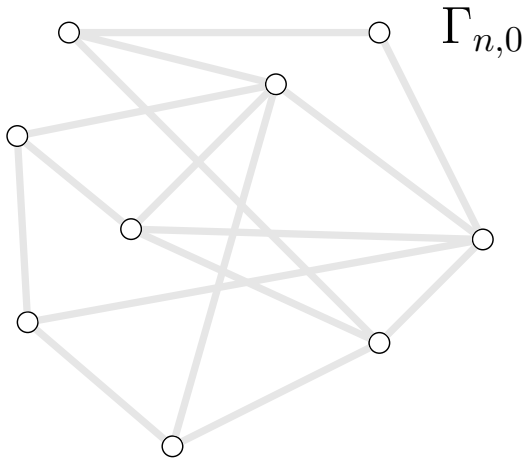
## The Erdős-Rényi process

- Let  $\Gamma = (V, E)$  be some **host graph** (from a sequence  $(\Gamma_n)$ )
- Suppose  $|V| = n$  and  $|E| = m$
- Order the edges  $E = (e_1, \dots, e_m)$  uniformly at random
- Define the **Erdős-Rényi process (with host  $\Gamma$ )** by

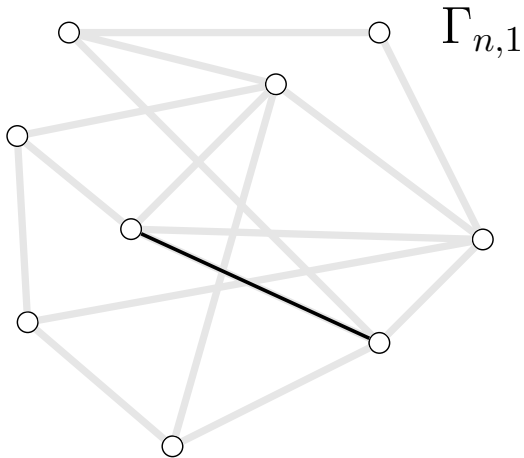
$$\Gamma_{n,t} = (V, \{e_1, e_2, \dots, e_t\}), \quad 0 \leq t \leq m$$

- Also define  $\Gamma_{n,p}$  (the **Erdős-Rényi subgraph**) as the graph obtained by deleting any edge from  $\Gamma$  independently with probability  $1 - p$
- $\Gamma_{n,t}$  and  $\Gamma_{n,p}$  with  $p = t/m$  are closely related

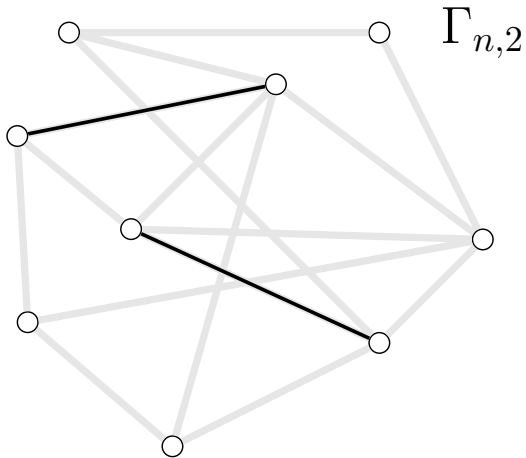
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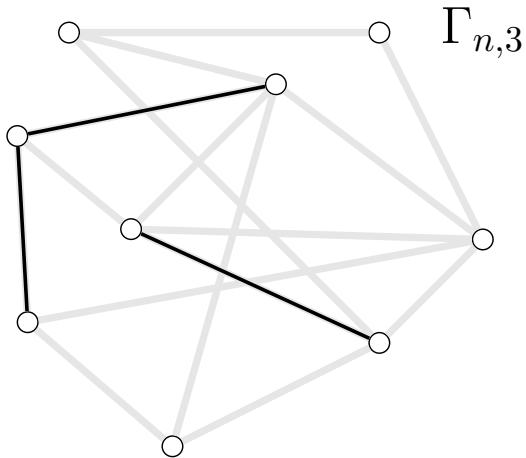
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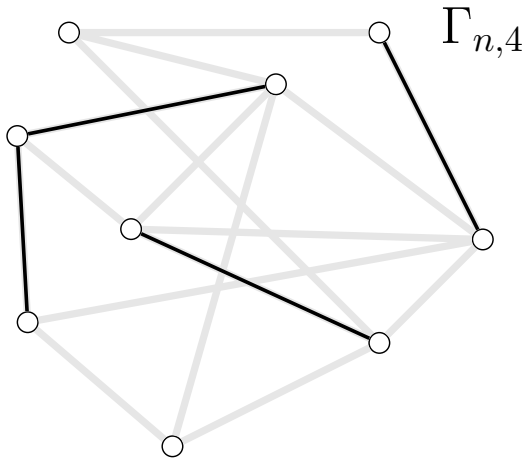


# The Erdős-Rényi process

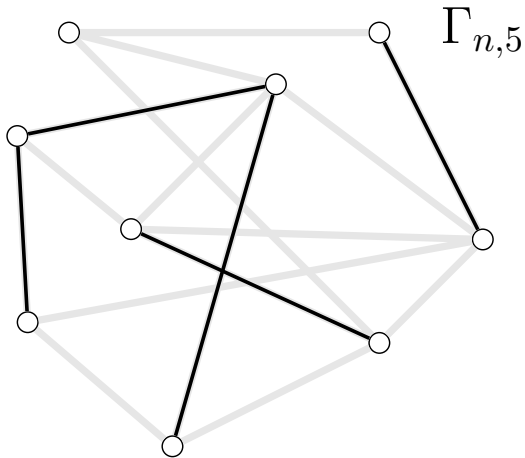




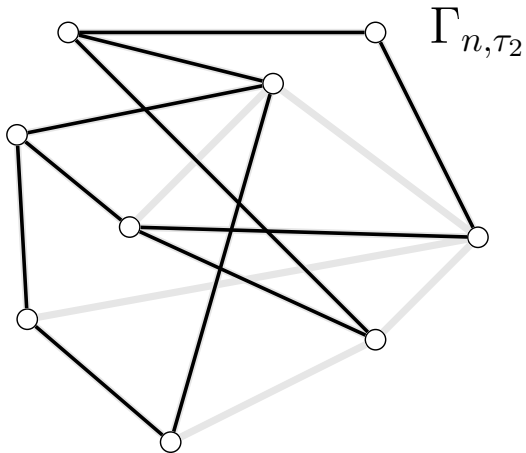
# The Erdős-Rényi process



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# The Erdős-Rényi process



## Hitting times

Given a Hamiltonian host  $\Gamma$ , define

$$\tau_2 = \min\{t : \delta(\Gamma_{n,t}) \geq 2\},$$

$$\tau_H = \min\{t : \Gamma_{n,t} \text{ Hamiltonian}\}.$$

**Theorem (Ajtai-Komlós-Szemerédi 1985)**

*If  $\Gamma = K_n$  then with high probability,*

$$\tau_2 = \tau_H.$$

In other words,  $G_{n,\tau_2}$  is Hamiltonian whp

## The threshold

Theorem (Pósa '76, Koršunov '76, Komlós-Szemerédi '83)

Suppose

$$p = \frac{\log n + \log \log n + c_n}{n}.$$

Then

$$\lim_{n \rightarrow \infty} \Pr \{G_{n,p} \text{ Hamiltonian}\} = \begin{cases} 0, & c_n \rightarrow -\infty, \\ e^{-e^{-c}}, & c_n \rightarrow c, \\ 1, & c_n \rightarrow \infty. \end{cases}$$

Frieze ('85) found a similar result for  $\Gamma = K_{n,n}$ .

## New results

Say  $(\Gamma_n)$  is a **Strong Dirac Graph sequence (SDG)**

if there exists some  $\varepsilon > 0$  such that  $\delta(\Gamma_n) \geq (1/2 + \varepsilon)n$  for all large  $n$ .

**Theorem (J. 2019+)**

*Suppose  $(\Gamma_n)$  is an SDG. Then with high probability,*

$$\tau_2 = \tau_H.$$

*The threshold for Hamiltonicity in  $\Gamma_{n,p}$  is the unique solution  $p_0$  to*

$$\sum_{v \in \Gamma_n} (1 - p)^{d_{\Gamma}(v)} \log n = 1. \quad (1)$$

LHS in (1)  $\approx$  expected number of vertices of degree less than 2.

## Regular host graphs

### Corollary

Suppose  $\Gamma_n$  is  $\beta n$ -regular for some  $\beta > 1/2$  (for all  $n$ ). Let

$$p = \frac{\log n + \log \log n + c_n}{\beta n}.$$

Then

$$\lim_{n \rightarrow \infty} \Pr \{ \Gamma_{n,p} \text{ Hamiltonian} \} = \begin{cases} 0, & c_n \rightarrow -\infty, \\ e^{-e^{-c}}, & c_n \rightarrow c, \\ 1, & c_n \rightarrow \infty. \end{cases}$$

“It is very pleasing, though perhaps not surprising, that a result similar to [P, K, KS] can be proved” – Frieze, '85

## What we will show

### Theorem

Suppose  $\omega \rightarrow \infty$  arbitrarily slowly with  $n$  and let

$$t = \tau_2 + \omega n.$$

Then

$$\lim_{n \rightarrow \infty} \Pr \{ \Gamma_{n,t} \text{ Hamiltonian} \} = 1.$$



## Proof plan

Will show:  $\Gamma_{n,t}$  Hamiltonian whp for  $t = \tau_2 + \omega n$

- Pósa's rotation lemma
- Pósa extensions
- Proof for  $\Gamma = K_n$
- Rotation-extension with smaller  $\Gamma$
- A colouring argument
- A more refined colouring argument

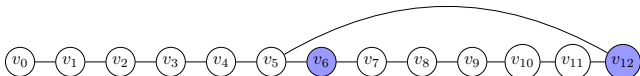
## Pósa's lemma

- Let  $G$  be a non-Hamiltonian graph
- Suppose  $P = (v_0, v_1, \dots, v_\ell)$  is a longest path in  $G$
- If  $v_\ell \sim v_i, i < \ell - 1$ , we obtain the path

$$P' = (v_0, v_1, \dots, v_i, v_\ell, v_{\ell-1}, \dots, v_{i+1})$$

via a **rotation with  $v_0$  fixed**

- $P'$  is also a longest path
- Let  $\text{EP}(v_0)$  be the set of endpoints obtainable with  $v_0$  fixed (so  $v_\ell, v_{i+1} \in \text{EP}(v_0)$ )



## Pósa's lemma

Lemma (Pósa '76)

$$|N(\text{EP}(v_0))| < 2|\text{EP}(v_0)|.$$

### Definition

A graph  $G$  **expands** if there is some  $\alpha > 0$  such that

$$|S| \leq \alpha n \quad \text{implies} \quad |N(S)| \geq 2|S|.$$

### Corollary

*In a non-Hamiltonian expander,  $|N(\text{EP}(v_0))| \geq \alpha n$ .*

## Proof plan

Will show:  $\Gamma_{n,t}$  Hamiltonian whp for  $t = \tau_2 + \omega n$

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## Pósa extension

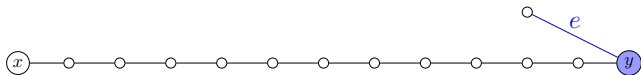
- Suppose  $G$  is a non-Hamiltonian, **connected** graph
- Suppose the longest path has length  $\ell$
- Let  $EP$  be the set of endpoints of  $\ell$ -length paths in  $G$

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- Suppose  $G$  is a non-Hamiltonian, **connected** graph
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- Add an edge  $e$  to  $G$ .

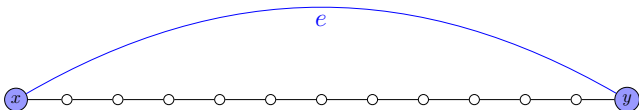
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## Pósa extension

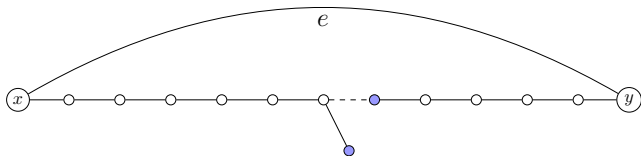
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- Add an edge  $e$  to  $G$ .
  - 1 If  $e$  connects  $EP$  to  $V \setminus P$ : path of length  $> \ell$
  - 2 If  $e$  connects some  $x \in EP$  to  $y \in EP(x)$ : cycle of length  $\ell + 1$ 
    - If  $\ell = n - 1$ , this is a Hamilton cycle.





## Pósa extension

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    - Otherwise, as  $G$  is connected, there exists a path of length  $> \ell$ .



## Pósa extension

- Suppose  $G$  is a non-Hamiltonian, **connected** graph
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  - ① If  $e$  connects  $EP$  to  $V \setminus P$ : path of length  $> \ell$
  - ② If  $e$  connects some  $x \in EP$  to  $y \in EP(x)$ : cycle of length  $\ell + 1$ 
    - If  $\ell = n - 1$ , this is a Hamilton cycle.
    - Otherwise, as  $G$  is connected, there exists a path of length  $> \ell$ .
  - ③ In other cases,  $e$  is not useful.

Say that  $e$  is an **extender** in cases 1 and 2.

## Proof plan

Will show:  $\Gamma_{n,t}$  Hamiltonian whp for  $t = \tau_2 + \omega n$

- Pósa's rotation lemma
  - Non-Hamiltonian expanders have  $|\text{EP}(x)| \geq \alpha n$
- Pósa extensions
  - Edge between  $x$  and  $\text{EP}(x)$  extends path (if  $G$  connected)
- **Proof for  $\Gamma = K_n$**
- Rotation-extension with smaller  $\Gamma$
- A colouring argument
- A more refined colouring argument

## Proof for $\Gamma = K_n$

- Easy:  $G_{n,t}$  is a connected expander whp for  $t = \frac{1}{2}(n \log n + n \log \log n + \omega n)$
- If  $G_{n,t}$  not already Hamiltonian, let  $x$  be an endpoint
- Consider the edge set

$$F = \bigcup_{y \in \text{EP}(x)} \{y\} \times \text{EP}(y).$$

We have at least  $|F| \geq \left(\frac{\alpha n}{2}\right)^2$  extenders

- Add a random edge  $e$  to form  $G_{n,t+1}$ .  
With probability  $\geq \alpha^2$ ,  $e$  is an extender.
- Whp, adding  $\omega n$  edges gives  $n$  extenders, so

$$\tau_H \leq t + \omega n$$

## Proof plan

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## Smaller host graphs

- Easy:  $\Gamma_{n,t}$  is a connected expander whp for  $t = p_0 m + \omega n$
- If  $\Gamma_{n,t}$  not already Hamiltonian, let  $x$  be an endpoint
- Consider the edge set

$$F = E(\Gamma) \cap \bigcup_{y \in \text{EP}(x)} \{y\} \times \text{EP}(y).$$

We might have  $|F| = 0$

The issue: we can only add edges from the host  $\Gamma$ ,  
and it may be that no extenders are in  $\Gamma$

**Example:**  $\Gamma = K_{n,n}$  and the longest path has odd length

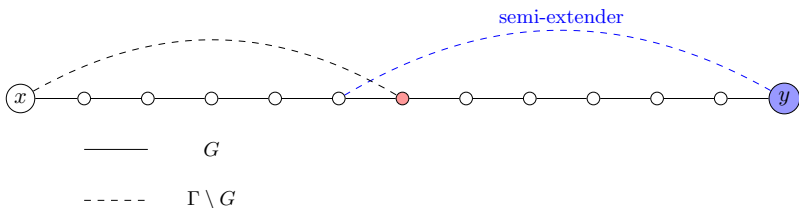
# Proof plan

Will show:  $\Gamma_{n,t}$  Hamiltonian whp for  $t = \tau_2 + \omega n$

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  - $\Omega(n^2)$  extenders  $\Rightarrow$  need  $\omega n$  edges for HC
- **Rotation-extension with smaller  $\Gamma$** 
  - $\Gamma$  may have  $o(n^2)$  extenders
- A colouring argument
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## Smaller host graphs

- Let  $G \subseteq \Gamma$  be a connected, non-Hamiltonian expander
- Suppose  $G$  has  $o(n^2)$  extenders with respect to  $\Gamma$
- Recall  $|\text{EP}(x)| \geq \alpha n$  for all  $x \in \text{EP}$
- Say  $x$  is **bad** if it has  $o(n)$  edges to  $\text{EP}(x)$  (in  $\Gamma$ )
- If  $x$  is bad and  $y \in \text{EP}(x)$ , this picture happens  $\geq \varepsilon n$  times:





## Creating extenders

- Let  $H(x)$  be the set of edges in  $\Gamma$  such that adding  $e \in H(x)$  increases the number of edges between  $x$  and  $\text{EP}(x)$
- As long as  $x$  is bad, this set has size  $\Omega(n^2)$
- Add  $\omega n$  random edges from  $\Gamma$  to  $G$
- If  $x$  is still bad, only  $o(n)$  edges from  $H(x)$  were added, but

$$\Pr \{ \text{Bin}(\omega n, 1 - \Omega(1)) = o(n) \} = e^{-\Omega(\omega n)}$$

- So with probability  $1 - e^{-\Omega(\omega n)}$ ,  $x$  is now good
- With probability  $1 - ne^{-\Omega(\omega n)} = 1 - e^{-\Omega(\omega n)}$ ,  
all  $x \in \text{EP}$  are now good

## Creating extenders

- We started with a connected non-Hamiltonian expander  $G \subseteq \Gamma$
- We added  $\omega n$  random edges from  $\Gamma$  to  $G$
- We now have  $\Omega(n^2)$  extenders wrt  $\Gamma$  in  $G$
- Adding another  $\omega$  edges will extend the path

**Issue:** This technique might require adding  $\omega n^2$  edges (impossible)

The argument will turn out to be useful anyway

## Smaller host graphs

Let's summarize the last few slides in a lemma.

We let  $\ell(G)$  denote the length of the longest path in  $G$ ,  
with  $\ell(G) = n$  if  $G$  is Hamiltonian.

### Lemma (“Extension lemma”)

*Suppose  $G \subseteq \Gamma$  is a connected non-Hamiltonian expander. Let  $G^+$  be the random graph obtained by adding  $\omega n$  random edges from  $\Gamma$  to  $G$ . Then*

$$\Pr \{ \ell(G^+) > \ell(G) \} = 1 - e^{-\Omega(\omega n)}$$

# Proof plan

Will show:  $\Gamma_{n,t}$  Hamiltonian whp for  $t = \tau_2 + \omega n$

- Pósa's rotation lemma
  - Non-Hamiltonian expanders have  $|\text{EP}(x)| \geq \alpha n$
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- **A colouring argument**
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## Observation

- **Extension lemma:** Adding  $\omega n$  random edges extends path with probability  $\geq 1 - e^{-\Omega(\omega n)}$
- **Reverse:** Independently delete edges with probability  $q = \omega / \log n$  (so about  $\omega n$  edges)
- Longest path survives with probability at least

$$(1 - q)^n = e^{-\Theta(\omega n / \log n)}$$

- The probabilities don't match, and we can use this

## A colouring argument

- We will define a random subgraph  $\Gamma_{n,t}^* \subseteq \Gamma_{n,t}$
- Define the event

$$\mathcal{A} = \{\ell(\Gamma_{n,t}^*) = \ell(\Gamma_{n,t}) < n\},$$

i.e. “neither graph is Hamiltonian and they share a longest path”

- Let

$$\mathcal{P}_k = \{\ell(\Gamma_{n,t}) = k\}, \quad k < n.$$

- Then

$$\Pr\{\mathcal{P}_k\} = \frac{\Pr\{\mathcal{A} \cap \mathcal{P}_k\}}{\Pr\{\mathcal{A} \mid \mathcal{P}_k\}}$$

## Choosing $\Gamma_{n,t}^*$

- Let  $q = \omega / \log n$
- Independently colour each edge in  $\Gamma_{n,t}$  **red** with probability  $q$ , and **blue** otherwise
- Let  $\Gamma_{n,t}^*$  be the **blue graph**, i.e. remove all red edges
- For any graph  $G$ ,

$$\Pr \{ \mathcal{A} \mid \Gamma_{n,t} = G \} \geq (1 - q)^k = e^{-\Theta(\omega n / \log n)}$$

(this is the probability that any fixed path of length  $k$  is all-**blue**)

- So for any event of the form  $\mathcal{E} = \{ \Gamma_{n,t} \in \mathcal{G} \}$ ,

$$\Pr \{ \mathcal{A} \mid \mathcal{E} \} \geq e^{-\Theta(\omega n / \log n)}$$

## The main calculation

With such a choice of  $\Gamma_{n,t}^*$  we have

$$\Pr \{ \mathcal{P}_k \} = \frac{\Pr \{ \mathcal{A} \cap \mathcal{P}_k \}}{\Pr \{ \mathcal{A} \mid \mathcal{P}_k \}} \leq \Pr \{ \mathcal{A} \cap \mathcal{P}_k \} e^{\Theta(\omega n / \log n)}$$

- We removed about  $qn \log n = \omega n$  red edges to form  $\Gamma_{n,t}^b$
- Conditioning on  $\Gamma_{n,t}^*$ , these are random and we may be able to use the extension lemma
- The extension lemma states that

$$\Pr \{ \mathcal{A} \mid \Gamma_{n,t}^* \text{ connected, non-Hamiltonian expander} \} \leq e^{-\Omega(\omega n)}$$

- **Aim:** use the lemma to show that

$$\Pr \{ \mathcal{A} \cap \mathcal{P}_k \} \leq e^{-\Omega(\omega n)}$$



## The main calculation

Define

$$\mathcal{P}_k^* = \{\ell(\Gamma_{n,t}^*) = k\}.$$

We have

$$\Pr\{\mathcal{A} \cap \mathcal{P}_k\} = \Pr\{\mathcal{P}_k^* \cap \mathcal{P}_k\} = \Pr\{\mathcal{P}_k \mid \mathcal{P}_k^*\} \Pr\{\mathcal{P}_k^*\} \leq \Pr\{\mathcal{A} \mid \mathcal{P}_k^*\}$$

But we **can't apply** extension lemma conditioning on  $\mathcal{P}_k^*$  only

Also need

$$\mathcal{C}^* = \{\Gamma_{n,t}^* \text{ is connected}\}$$

$$\mathcal{E}^* = \{\Gamma_{n,t}^* \text{ expands}\}$$

# Proof plan

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- A colouring argument
  - Need  $\Gamma_{n,t}^*$  to be a connected expander
- **A more refined colouring argument**

## Bringing in expansion

More refined approach: for any events  $\mathcal{M}, \mathcal{N}$ ,

$$\Pr\{\mathcal{P}_k \cap \mathcal{N}\} = \frac{\Pr\{\mathcal{A} \cap \mathcal{M} \cap \mathcal{P}_k \cap \mathcal{N}\}}{\Pr\{\mathcal{A} \cap \mathcal{M} \mid \mathcal{P}_k \cap \mathcal{N}\}}$$

We will carefully pick  $\mathcal{M}, \mathcal{N}$ .

Plan is to apply

- extension lemma to numerator,
- $\Pr\{\mathcal{A} \mid \Gamma_{n,t}\} \geq (1 - q)^n$  to denominator
- Error terms will appear

## Picking $\mathcal{M}, \mathcal{N}$

$$\Pr \{ \mathcal{P}_k \cap \mathcal{N} \} = \frac{\Pr \{ \mathcal{A} \cap \mathcal{M} \cap \mathcal{P}_k \cap \mathcal{N} \}}{\Pr \{ \mathcal{A} \cap \mathcal{M} \mid \mathcal{P}_k \cap \mathcal{N} \}}$$

What we need for the **denominator**:

- $\mathcal{N}$  should only concern  $\Gamma_{n,t}$  so that we have

$$\Pr \{ \mathcal{A} \mid \mathcal{P}_k \cap \mathcal{N} \} \geq (1 - q)^n = e^{-\Theta(\omega n / \log n)}$$

- We will use

$$\Pr \{ \mathcal{A} \cap \mathcal{M} \mid \mathcal{P}_k \cap \mathcal{N} \} \geq \Pr \{ \mathcal{A} \mid \mathcal{P}_k \cap \mathcal{N} \} - \Pr \{ \overline{\mathcal{M}} \mid \mathcal{P}_k \cap \mathcal{N} \},$$

so we need

$$\Pr \{ \overline{\mathcal{M}} \mid \mathcal{P}_k \cap \mathcal{N} \} \leq e^{-\Omega(\omega n / \log n)}$$

## Picking $\mathcal{M}, \mathcal{N}$

$$\Pr \{ \mathcal{P}_k \cap \mathcal{N} \} = \frac{\Pr \{ \mathcal{A} \cap \mathcal{M} \cap \mathcal{P}_k \cap \mathcal{N} \}}{\Pr \{ \mathcal{A} \cap \mathcal{M} \mid \mathcal{P}_k \cap \mathcal{N} \}}$$

What we need for the **numerator**:

- Let  $\mathcal{B}_k = (\mathcal{M} \cap \mathcal{N} \cap \mathcal{P}_k) \setminus (\mathcal{C}^* \cap \mathcal{E}^*)$ . Then

$$\Pr \{ \mathcal{A} \cap \mathcal{M} \cap \mathcal{P}_k \cap \mathcal{N} \} \leq \Pr \{ \mathcal{A} \cap \mathcal{P}_k \cap \mathcal{C}^* \cap \mathcal{E}^* \} + \Pr \{ \mathcal{B}_k \}$$

- First term, using the extension lemma, is bounded by

$$\Pr \{ \mathcal{A} \mid \mathcal{C}^* \cap \mathcal{E}^* \cap \mathcal{P}_k^* \} \leq e^{-\Omega(\omega n)}$$

- Need

$$\Pr \{ \mathcal{B}_k \} \leq e^{-\Omega(\omega n)}$$

## Constraint roundup

We also need  $\Pr \{\mathcal{N}\} = 1 - o(1)$  so that

$$\Pr \{\Gamma_{n,t} \text{ not Hamiltonian}\} \leq \Pr \{\overline{\mathcal{N}}\} + \sum_k \Pr \{\mathcal{P}_k \cap \mathcal{N}\} = o(1).$$

Pick events  $\mathcal{M}, \mathcal{N}$  such that

- $\mathcal{N}$  concerns only  $\Gamma_{n,t}$ ,
- $\Pr \{\mathcal{N}\} = 1 - o(1)$ ,
- $\Pr \{\overline{\mathcal{M}} \mid \mathcal{P}_k \cap \mathcal{N}\} \leq e^{-\Omega(\omega n / \log n)}$ ,
- $\Pr \{(\mathcal{M} \cap \mathcal{N} \cap \mathcal{P}_k) \setminus (\mathcal{C}^* \cap \mathcal{E}^*)\} \leq e^{-\Omega(\omega n)}$ .

## The underlying challenge

- For  $t$  in our range, we only have

$$\Pr \{ \Gamma_{n,t} \text{ expands} \} = 1 - o(n^{-1+\varepsilon})$$

and even worse for  $\Gamma_{n,t}^*$

- We want to say something like

$$\Pr \{ \Gamma_{n,t}^* \text{ expands} \mid \mathcal{P}_k \cap \mathcal{N} \} = 1 - e^{-\Omega(\omega n)}$$

- If expansion were a decreasing property we could take

$$\mathcal{N} = \{ \Gamma_{n,t} \text{ expands} \}$$

- We will redefine  $\Gamma_{n,t}^*$  so that it expands when  $\Gamma_{n,t}$  does

## Redefine $\Gamma_{n,t}^*$

- Colour edges in  $\Gamma_{n,t}$  **red** with probability  $q = \frac{\omega}{\log n}$ , else **blue**
- Let

$$\text{SMALL} = \left\{ v : \deg(v) \leq \frac{\log n}{100} \text{ in blue subgraph} \right\}$$

- Let  $\text{LARGE} = V \setminus \text{SMALL}$
- Remove all red edges fully contained in LARGE
- The resulting graph is  $\Gamma_{n,t}^*$
- $\Gamma_{n,t}^*$  and  $\Gamma_{n,t}$  agree on properties only concerning SMALL and edges incident to them



## On expansion

### Lemma

We have

$$\mathcal{D} \cap \mathcal{E}_1 \cap \mathcal{E}_2 \subseteq \{\Gamma_{n,t} \text{ expands}\}$$

where

$$\mathcal{D} = \{\delta(\Gamma_{n,t}) \geq 2\},$$

$$\mathcal{E}_1 = \{\text{every vertex set } |S| \leq 6\alpha n \text{ has } e(S) \leq \frac{\log n}{1000} |S| \text{ in } \Gamma_{n,t}\},$$

$$\mathcal{E}_2 = \{\Gamma_{n,t} \text{ contains no paths or cycles of length } \leq 4 \text{ involving two edges incident to SMALL}\}.$$

$$(\alpha = e^{-2000})$$

**The point:**  $\mathcal{D}, \mathcal{E}_1, \mathcal{E}_2$  immediately transfer to  $\Gamma_{n,t}^*$

## Picking $\mathcal{M}, \mathcal{N}$

We let

$$\mathcal{N} = \{\delta(\Gamma_{n,t}) \geq 2\} \cap \mathcal{E}_1 \cap \mathcal{E}_2$$

$$\mathcal{M} = \{\Gamma_{n,t}^* \text{ connected}\}$$

What we needed was:

- $\mathcal{N}$  concerns only  $\Gamma_{n,t}$  (by design)
- $\Pr\{\mathcal{N}\} = 1 - o(1)$  (routine bounds)
- $\Pr\{\overline{\mathcal{M}} \mid \mathcal{P}_k \cap \mathcal{N}\} \leq e^{-\Omega(\omega n / \log n)}$   
(expansion makes connectivity likely)
- $\Pr\{(\mathcal{M} \cap \mathcal{N} \cap \mathcal{P}_k) \setminus (\mathcal{C}^* \cap \mathcal{E}^*)\} \leq e^{-\Omega(\omega n)}$   
(is actually zero since  $\mathcal{M} = \mathcal{C}^*$ ,  $\mathcal{N} \subseteq \mathcal{E}^*$ )

## We are done

We have

$$\Pr \{ \Gamma_{n,t} \text{ not Hamiltonian} \} \leq \Pr \{ \overline{\mathcal{N}} \} + \sum_{k < n} \Pr \{ \mathcal{P}_k \cap \mathcal{N} \}$$

and

$$\begin{aligned} \Pr \{ \mathcal{P}_k \cap \mathcal{N} \} &= \frac{\Pr \{ \mathcal{A} \cap \mathcal{M} \cap \mathcal{P}_k \cap \mathcal{N} \}}{\Pr \{ \mathcal{A} \cap \mathcal{M} \mid \mathcal{P}_k \cap \mathcal{N} \}} \\ &\leq \frac{e^{-\Omega(\omega n)}}{e^{-\Theta(\omega n / \log n)}} \\ &= o(n^{-1}). \end{aligned}$$

# Proof plan

Will show:  $\Gamma_{n,t}$  Hamiltonian whp for  $t = \tau_2 + \omega n$

- Pósa's rotation lemma
  - Non-Hamiltonian expanders have  $|\text{EP}(x)| \geq \alpha n$
- Pósa extensions
  - Edge between  $x$  and  $\text{EP}(x)$  extends path (if  $G$  connected)
- Proof for  $\Gamma = K_n$ 
  - $\Omega(n^2)$  extenders  $\Rightarrow$  need  $\omega n$  edges for HC
- Rotation-extension with smaller  $\Gamma$ 
  - $\Gamma$  may have  $o(n^2)$  extenders
  - Extension lemma: need  $\omega n$  edges to get longer path
- A colouring argument
  - Need  $\Gamma_{n,t}^*$  to be a connected expander
- A more refined colouring argument
  - Make  $\Gamma_{n,t}^*$ ,  $\Gamma_{n,t}$  agree on low-degree vertices
  - Make expansion transfer from  $\Gamma_{n,t}$  to  $\Gamma_{n,t}^*$

## Summary

Why does this work?

- Most edges are between vertices of high degree ( $\Omega(\log n)$ )
- Such edges are unlikely to be in longest path
- Randomly removing  $\omega n$  such edges is fairly likely to be safe (longest path remains and the graph still expands)
- Randomly throwing them back in is likely to extend path
- This “contradiction” is formalized by

$$\Pr \{ \mathcal{P}_k \cap \mathcal{N} \} = \frac{\Pr \{ \mathcal{A} \cap \mathcal{M} \cap \mathcal{P}_k \cap \mathcal{N} \}}{\Pr \{ \mathcal{A} \cap \mathcal{M} \mid \mathcal{P}_k \cap \mathcal{N} \}}$$

- This argument seems suited for **spanning structures** in **random subgraphs** of **non-complete host graphs**

## Summary

What were the keys?

- We redefined  $\Gamma_{n,t}^*$  so that it agrees with  $\Gamma_{n,t}$  on vertices of low degree
- We considered a stricter type of expansion which automatically transfers from  $\Gamma_{n,t}$  to  $\Gamma_{n,t}^*$

Things I skipped:

- How expansion follows from decreasing events and  $\delta \geq 2$
- We can consider  $\Gamma_{n,\tau_2}$  without too much trouble, but it adds even more events to the calculations
- We should have conditioned on  $\Omega(\omega n)$  red edges being removed
- When we re-place the red edges, they must be in LARGE.  
Should have argued that  $|\text{LARGE}| = n - o(n)$  (more events)